

# Recent Advance in Numerical Algorithms for Large/Sparse Eigenvalue Problems

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# 前言

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- 1 Power Method
- 2 Inverse Power Iteration

# 前言

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## Eigenproblems

- Standard eigenproblems:  $Ax = \lambda x$
- Generalized eigenproblems:  $Ax = \lambda Bx$
- Higher order poly. eigenproblems:  
 $(A_0 + \lambda A_1 + \dots + \lambda^n A_n)x = 0$
- Eigenproblems of  $\lambda$ -matrices:  $F(\lambda)x = 0$

## What do we care ?

- (i) In theory: eigenstructure, spectral decomposition, canonical form, ..., etc.
- (ii) In computation: eigenvalues, eigenvectors, invariant subspaces, ..., etc.

# Power Methods

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Given  $A \in \mathbb{C}^{n \times n}$ . Let  $A$  be diagonalizable and

$$Ax_i = \lambda_i x_i, \quad i = 1, \dots, n$$

with

$$|\lambda_1| > |\lambda_2| \geq \dots \geq |\lambda_n|.$$

## Goal

Find the maximum eigenvalue  $\lambda_1$  and the corresponding eigenvector  $x_1$ .

# Power Methods

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Let  $u_0 \neq 0$  be a given vector. From the expansion

$$u_0 = \sum_{i=1}^n \alpha_i x_i, \quad \alpha_1 \neq 0,$$

follows that

$$\begin{aligned} A^k u_0 &= \sum_{i=1}^n \alpha_i \lambda_i^k x_i \\ &= \lambda_1^k \left\{ \alpha_1 x_1 + \sum_{i=2}^n \alpha_i \left( \frac{\lambda_i}{\lambda_1} \right)^k x_i \right\}. \end{aligned} \quad (1)$$

Thus

$$\lambda_1^{-k} A^k u_0 \rightarrow \alpha_1 x_1 \quad \text{as } k \rightarrow \infty.$$

# Power Methods

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## Algorithm (Power Method with 2-norm)

Choose an initial  $u \neq 0$  with  $\|u\|_2 = 1$ .

Iterate until convergence

Compute  $v = Au$ ;  $m = \|v\|_2$ ;  $u := v/m$

## Theorem

The sequence defined by Algorithm 1 is satisfied

$$\lim_{k \rightarrow \infty} m_k = |\lambda_1|,$$

$$\lim_{k \rightarrow \infty} \varepsilon^k u_k = \frac{x_1}{\|x_1\|} \frac{\alpha_1}{|\alpha_1|}, \quad \text{where } \varepsilon = \frac{|\lambda_1|}{\lambda_1}.$$

# Power Methods

*Proof.* It is obvious that

$$u_k = A^k u_0 / \|A^k u_0\|, \quad m_k = \|A^k u_0\| / \|A^{k-1} u_0\|. \quad (2)$$

This follows from  $\lambda_1^{-k} A^k u_0 \rightarrow \alpha_1 x_1$  that

$$\begin{aligned} |\lambda_1|^{-k} \|A^k u_0\| &\rightarrow |\alpha_1| \|x_1\| \\ |\lambda_1|^{-k+1} \|A^{k-1} u_0\| &\rightarrow |\alpha_1| \|x_1\| \end{aligned}$$

and then

$$|\lambda_1|^{-1} \|A^k u_0\| / \|A^{k-1} u_0\| = |\lambda_1|^{-1} m_k \rightarrow 1.$$

From (1) follows now for  $k \rightarrow \infty$

$$\begin{aligned} \varepsilon^k u_k &= \varepsilon^k \frac{A^k u_0}{\|A^k u_0\|} = \frac{\alpha_1 x_1 + \sum_{i=2}^n \alpha_i \left(\frac{\lambda_i}{\lambda_1}\right)^k x_i}{\|\alpha_1 x_1 + \sum_{i=2}^n \alpha_i \left(\frac{\lambda_i}{\lambda_1}\right)^k x_i\|} \\ &\rightarrow \frac{\alpha_1 x_1}{\|\alpha_1 x_1\|} = \frac{x_1}{\|x_1\|} \frac{\alpha_1}{|\alpha_1|}. \end{aligned}$$

# Power Methods

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## Algorithm (Power Method with Linear Function)

Choose an initial  $u_0 \neq 0$ .

Iterate until convergence

Compute  $v_{k+1} = Au_k$ ,  $m_{k+1} = \ell(v_{k+1})$ ,  $u_{k+1} = v_{k+1}/m_{k+1}$   
where  $\ell(v_{k+1})$ , e.g.  $e_1(v_{k+1})$  or  $e_n(v_{k+1})$ , is a linear functional.

## Theorem

Suppose  $\ell(x_1) \neq 0$  and  $\ell(v_k) \neq 0$ ,  $k = 1, 2, \dots$ , then

$$\lim_{k \rightarrow \infty} m_k = \lambda_1$$

$$\lim_{k \rightarrow \infty} u_k = \frac{x_1}{\ell(x_1)}.$$

# Power Methods

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*Proof.* As above we show that

$$u_k = A^k u_0 / \ell(A^k u_0), \quad m_k = \ell(A^k u_0) / \ell(A^{k-1} u_0).$$

From (1) we get for  $k \rightarrow \infty$

$$\lambda_1^{-k} \ell(A^k u_0) \rightarrow \alpha_1 \ell(x_1),$$

$$\lambda_1^{-k+1} \ell(A^{k-1} u_0) \rightarrow \alpha_1 \ell(x_1),$$

thus

$$\lambda_1^{-1} m_k \rightarrow 1, \quad k \rightarrow \infty.$$

Similarly for  $k \rightarrow \infty$ ,

$$u_k = \frac{A^k u_0}{\ell(A^k u_0)} = \frac{\alpha_1 x_1 + \sum_{i=2}^n \alpha_i \left(\frac{\lambda_i}{\lambda_1}\right)^k x_i}{\ell(\alpha_1 x_1 + \sum_{i=2}^n \alpha_i \left(\frac{\lambda_i}{\lambda_1}\right)^k x_i)} \rightarrow \frac{\alpha_1 x_1}{\alpha_1 \ell(x_1)}$$

# Power Methods

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- Note that:

$$\begin{aligned} m_k &= \frac{\ell(A^k u_0)}{\ell(A^{k-1} u_0)} = \lambda_1 \frac{\alpha_1 \ell(x_1) + \sum_{i=2}^n \alpha_i \left(\frac{\lambda_i}{\lambda_1}\right)^k \ell(x_i)}{\alpha_1 \ell(x_1) + \sum_{i=2}^n \alpha_i \left(\frac{\lambda_i}{\lambda_1}\right)^{k-1} \ell(x_i)} \\ &= \lambda_1 + O\left(\left|\frac{\lambda_2}{\lambda_1}\right|^{k-1}\right). \end{aligned}$$

That is the convergent rate is  $\left|\frac{\lambda_2}{\lambda_1}\right|$ .

# 二、Inverse Power Iteration

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## Goal

Find the eigenvalue of  $A$  that is in a given region or closest to a certain scalar  $\sigma$  and the corresponding eigenvector.

- This is achieved by iterating with the matrix  $(A - \sigma I)^{-1}$ .
- Often,  $\sigma$  is referred to as the “Shift”.

## Algorithm (Inverse power method with a fixed shift)

Choose an initial  $u_0 \neq 0$ .

For  $k = 0, 1, 2, \dots$

Compute  $v_{k+1} = (A - \sigma I)^{-1}u_k$  and  $m_{k+1} = \ell(v_{k+1})$ .

Set  $u_{k+1} = v_{k+1}/m_{k+1}$

# 二、Inverse Power Iteration

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- The convergence of Algorithm 3 is  $|\frac{\lambda_1 - \sigma}{\lambda_2 - \sigma}|$  whenever  $\lambda_1$  and  $\lambda_2$  are the closest and the second closest eigenvalues to  $\sigma$ .
- Algorithm 3 is linearly convergent.

## Algorithm (Inverse power method with variant shifts)

Choose an initial  $u_0 \neq 0$ .

Given  $\sigma_0 = \sigma$ .

For  $k = 0, 1, 2, \dots$

Compute  $v_{k+1} = (A - \sigma_k I)^{-1} u_k$  and  $m_{k+1} = \ell(v_{k+1})$ .

Set  $u_{k+1} = v_{k+1}/m_{k+1}$  and  $\sigma_{k+1} = \sigma_k + 1/m_{k+1}$ .

- Above algorithm is locally quadratic convergent.

# ≡ Connection with Newton method

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Consider the nonlinear equations:

$$F \left( \begin{bmatrix} u \\ \lambda \end{bmatrix} \right) \equiv \begin{bmatrix} Au - \lambda u \\ \ell^T u - 1 \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \end{bmatrix}. \quad (3)$$

Newton method for (3): for  $k = 0, 1, 2, \dots$

$$\begin{bmatrix} u_{k+1} \\ \lambda_{k+1} \end{bmatrix} = \begin{bmatrix} u_k \\ \lambda_k \end{bmatrix} - \left[ F' \left( \begin{bmatrix} u_k \\ \lambda_k \end{bmatrix} \right) \right]^{-1} F \left( \begin{bmatrix} u_k \\ \lambda_k \end{bmatrix} \right).$$

Since

$$F' \left( \begin{bmatrix} u \\ \lambda \end{bmatrix} \right) = \begin{bmatrix} A - \lambda I & -u \\ \ell^T & 0 \end{bmatrix},$$

the Newton method can be rewritten by component-wise

$$(A - \lambda_k)u_{k+1} = (\lambda_{k+1} - \lambda_k)u_k \quad (4)$$

$$\ell^T u_{k+1} = 1. \quad (5)$$

# ≡、Connection with Newton method

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Let

$$v_{k+1} = \frac{u_{k+1}}{\lambda_{k+1} - \lambda_k}.$$

Substituting  $v_{k+1}$  into (4), we get

$$(A - \lambda_k I)v_{k+1} = u_k.$$

By equation (5), we have

$$m_{k+1} = \ell(v_{k+1}) = \frac{\ell(u_{k+1})}{\lambda_{k+1} - \lambda_k} = \frac{1}{\lambda_{k+1} - \lambda_k}.$$

It follows that

$$\lambda_{k+1} = \lambda_k + \frac{1}{m_{k+1}}.$$

Hence the Newton's iterations (4) and (5) are identified with Algorithm 4.

# ≡、Connection with Newton method

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## Theorem

Let  $u \neq 0$  and for any  $\mu$  set  $r_\mu = Au - \mu u$ . Then  $\|r_\mu\|_2$  is minimized when

$$\mu = \theta = u^* Au / u^* u.$$

In this case  $r_\theta \perp u$ .

*Proof:* W.L.O.G. assume  $\|u\|_2 = 1$ . Let  $[u \ U]$  be unitary and set

$$\begin{bmatrix} u^* \\ U^* \end{bmatrix} A \begin{bmatrix} u & U \end{bmatrix} = \begin{bmatrix} u^* Au & u^* AU \\ U^* Au & U^* AU \end{bmatrix} \equiv \begin{bmatrix} \theta & h^* \\ g & B \end{bmatrix}.$$

# 三、Connection with Newton method

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Then

$$\begin{aligned} \begin{bmatrix} u^* \\ U^* \end{bmatrix} r_\mu &= \begin{bmatrix} u^* \\ U^* \end{bmatrix} Au - \mu \begin{bmatrix} u^* \\ U^* \end{bmatrix} u \\ &= \begin{bmatrix} u^* \\ U^* \end{bmatrix} A \begin{bmatrix} u & U \end{bmatrix} \begin{bmatrix} u^* \\ U^* \end{bmatrix} u - \mu \begin{bmatrix} u^* \\ U^* \end{bmatrix} u \\ &= \begin{bmatrix} \theta & h^* \\ g & B \end{bmatrix} \begin{bmatrix} u^* \\ U^* \end{bmatrix} u - \mu \begin{bmatrix} u^* \\ U^* \end{bmatrix} u \\ &= \begin{bmatrix} \theta & h^* \\ g & B \end{bmatrix} \begin{bmatrix} 1 \\ 0 \end{bmatrix} - \mu \begin{bmatrix} 1 \\ 0 \end{bmatrix} = \begin{bmatrix} \theta - \mu \\ g \end{bmatrix}. \end{aligned}$$

It implies that

$$\|r_\mu\|_2^2 = \left\| \begin{bmatrix} u^* \\ U^* \end{bmatrix} r_\mu \right\|_2^2 = \left\| \begin{bmatrix} \theta - \mu \\ g \end{bmatrix} \right\|_2^2 = |\theta - \mu|^2 + \|g\|_2^2.$$

Therefore

$$\min_{\mu} \|r_\mu\|_2 = \|g\|_2 = \|r_\theta\|_2.$$

# ≡ Connection with Newton method

That is

$$\mu = \theta = u^* Au,$$

and

$$u^* r_\theta = u^*(Au - \theta u) = u^* Au - \theta = 0.$$

Hence,  $r_\theta \perp u$ . □

## Definition

Let  $u$  be a nonzero vector. Then  $u^* Au / u^* u$  is called a **Rayleigh quotient**.

If  $u$  is an eigenvector corresponding to an eigenvalue  $\lambda$  of  $A$ , then

$$\frac{u^* Au}{u^* u} = \frac{\lambda u^* u}{u^* u} = \lambda.$$

$\Rightarrow u_k^* Au_k / u_k^* u_k$  provide a sequence of approximation to  $\lambda$  in the power method.

# ≡、Connection with Newton method

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## Algorithm (Inverse power method with Rayleigh Quotient)

Choose an initial  $u_0 \neq 0$  with  $\|u_0\|_2 = 1$ .

Compute  $\sigma_0 = u_0^\top A u_0$ .

For  $k = 0, 1, 2, \dots$

Compute  $v_{k+1} = (A - \sigma_k I)^{-1} u_k$ .

Set  $u_{k+1} = v_{k+1} / \|v_{k+1}\|_2$  and  $\sigma_{k+1} = u_{k+1}^\top A u_{k+1}$ .

- For symmetric  $A$ , above algorithm is cubically convergent.